

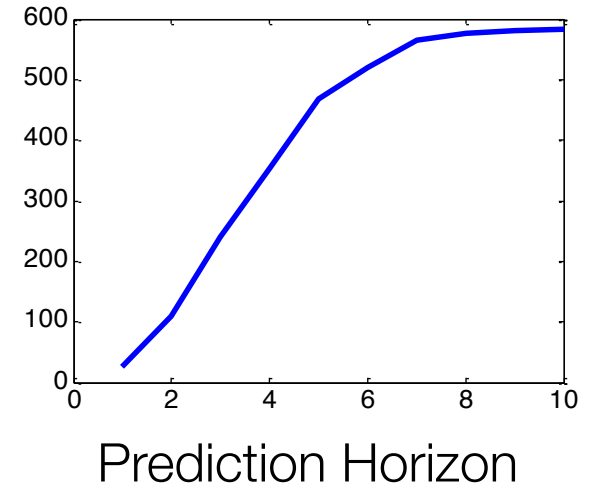
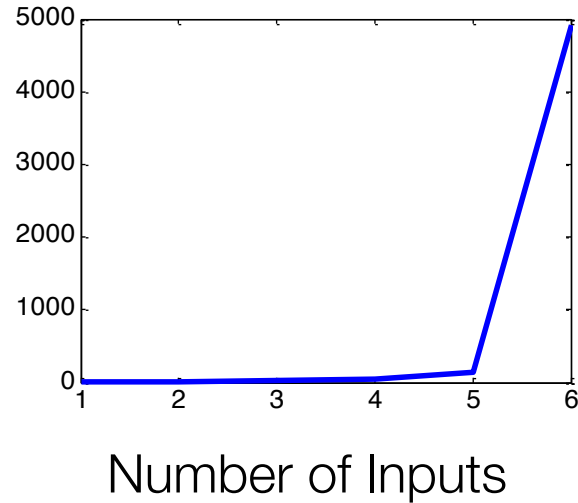
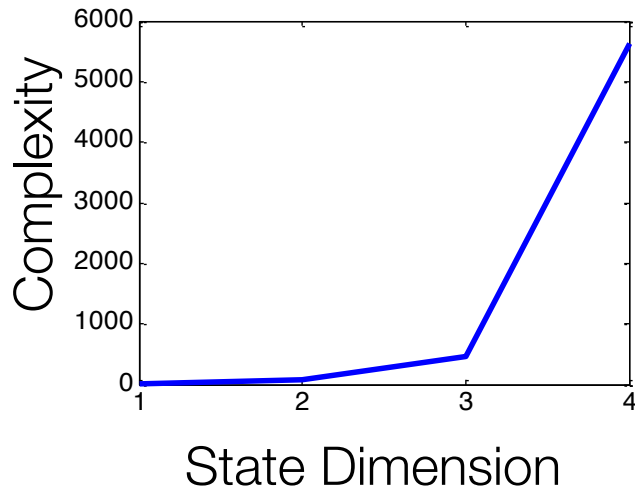
Fixed Complexity Explicit Model Predictive Control

Colin Jones and Melanie Zeilinger

 Automatic Control Laboratory, EPFL

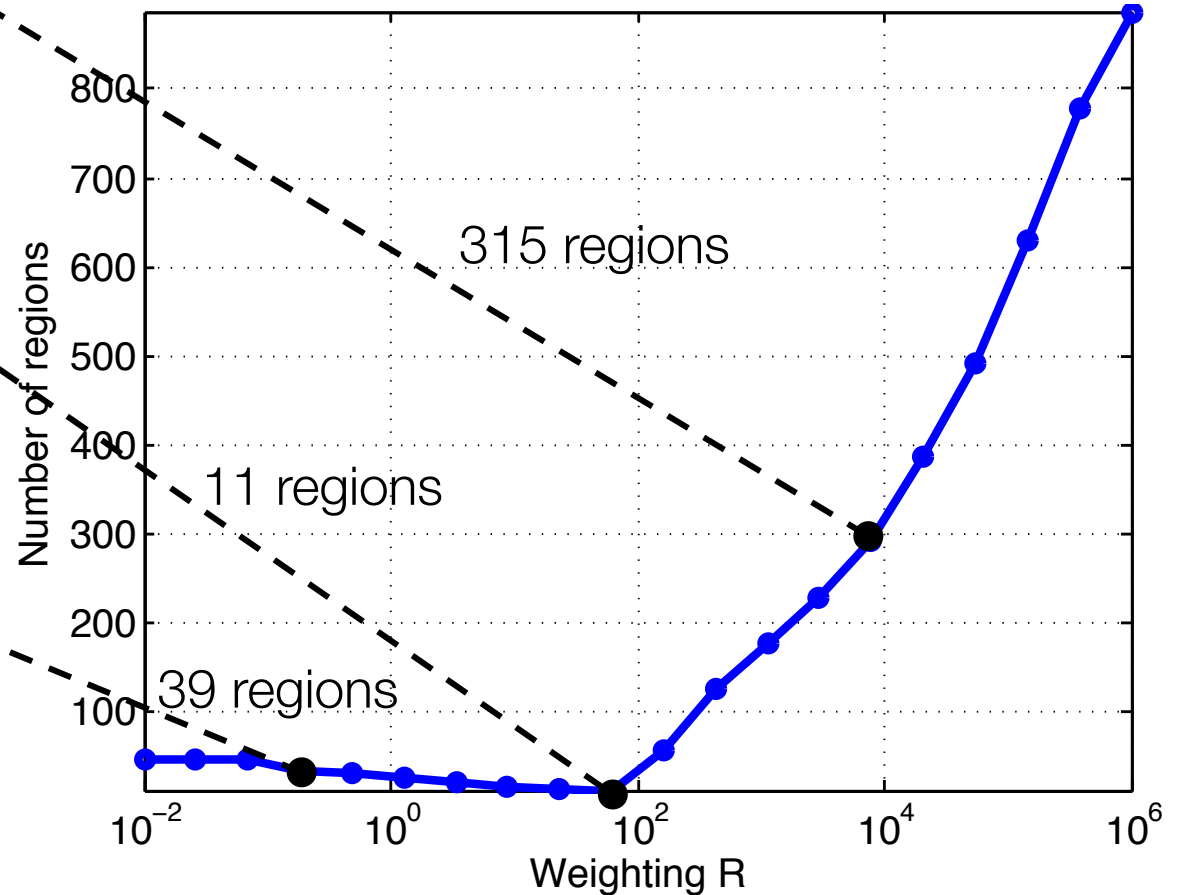
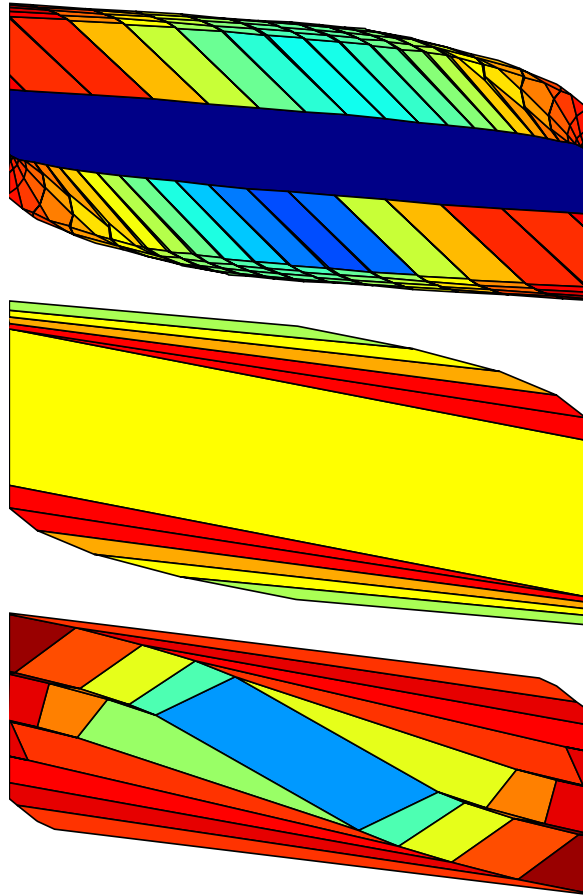
Complexity highly sensitive to problem size

Rapid increase in complexity with parameters



Impact of Tuning is Highly Uncertain

- Two-dimensional MPC problem with one input
- Tune the weighting matrix 'R' on the input



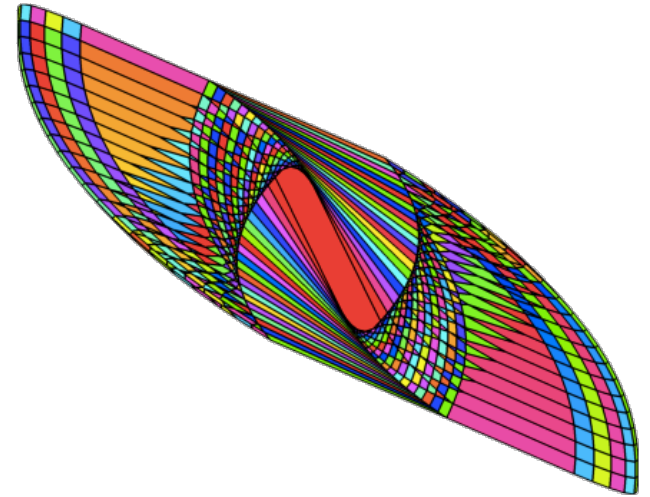
Limiting Factor: Complexity

Number of regions determine important properties

- online computation time
- storage requirements
- offline processing time

Complexity is a property of the problem

- Nothing is known about relationship between problem parameters and complexity

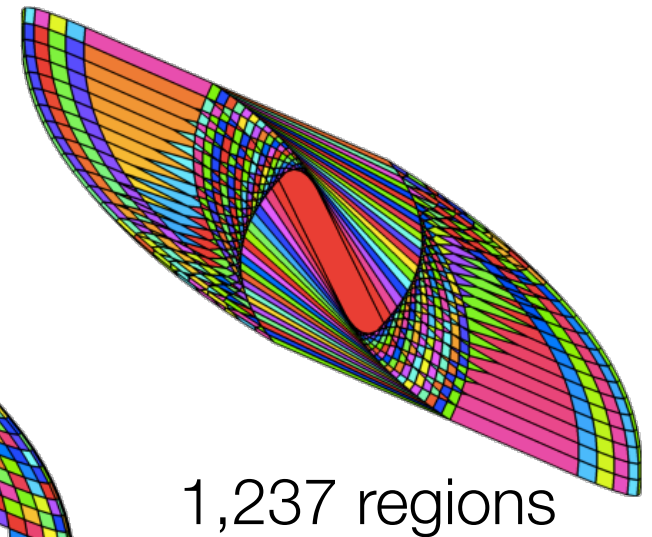
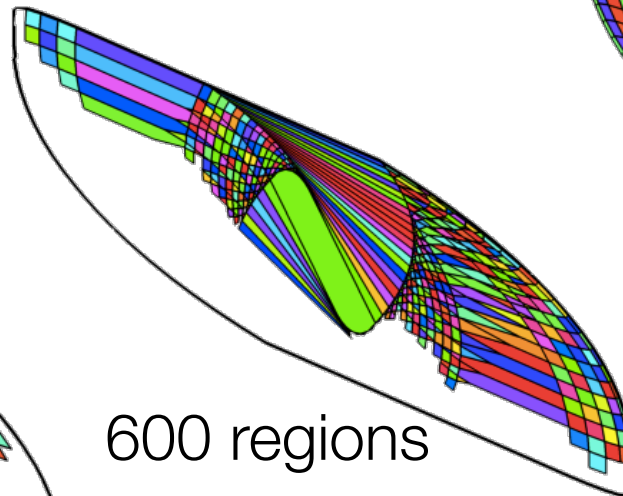
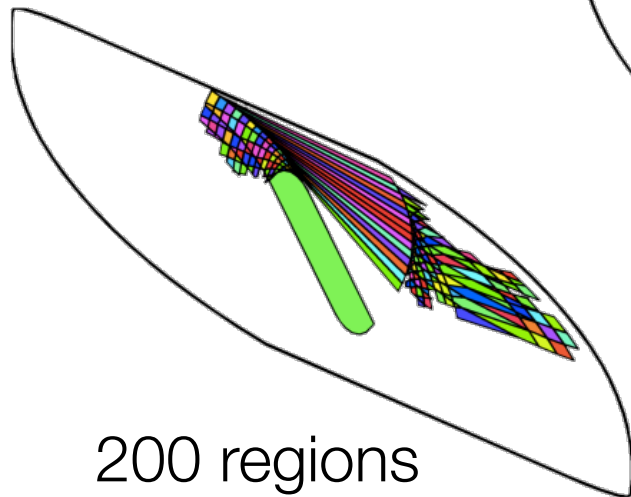


Fixed-complexity explicit MPC: Sub-optimal controller of exactly M regions

- Is it stable?
- Invariant?
- What level of sub-optimality?

Current optimal algorithms

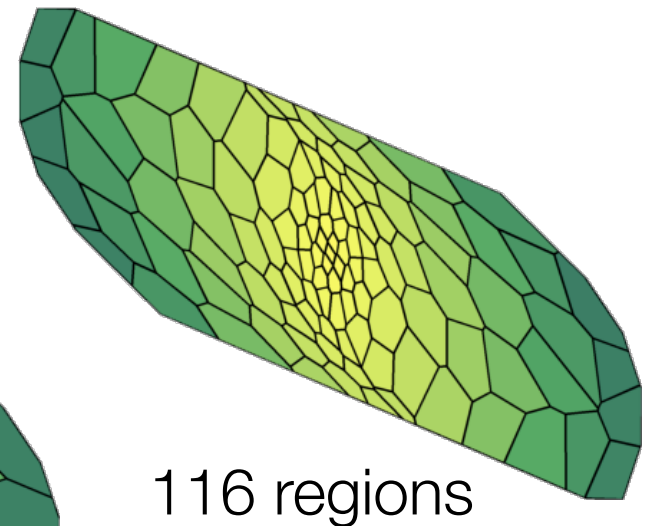
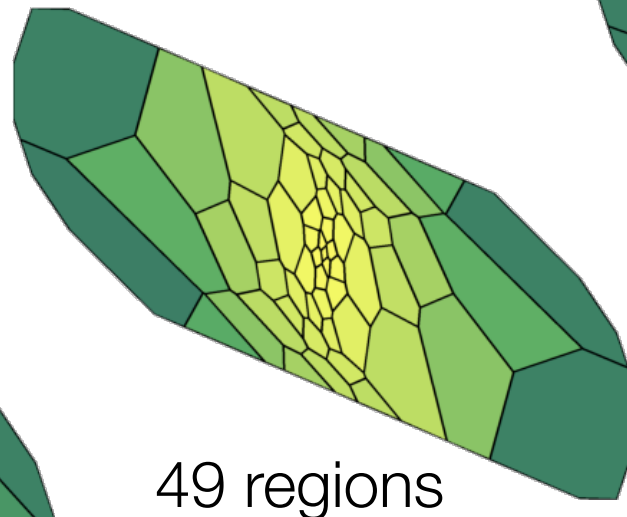
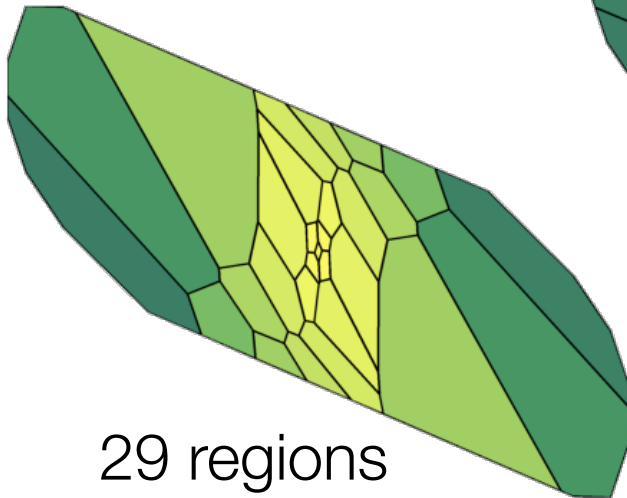
- Hardware can only store M regions
- Current algorithms are non-incremental
 - output meaningless until complete
 - cannot stop early: M regions gives nothing



Goal: Incremental algorithms

Assume: Embedded processor can store M regions

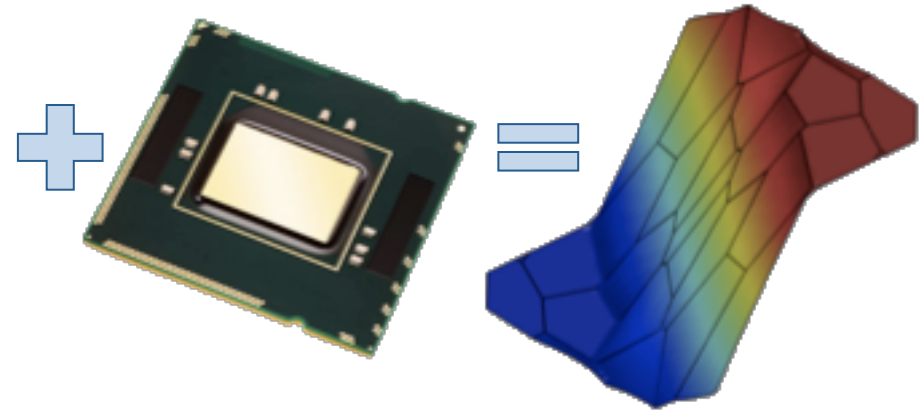
- Initially poor approximation
- Improve performance by adding complexity
- Stop when M regions reached
 - Certificate of stability/performance



Real-time synthesis : Complexity as a specification

$$u^*(x) = \underset{u_i}{\operatorname{argmin}} V_N(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$

s.t. $x_{i+1} = f(x_i, u_i)$
 $(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$
 $x_N \in \mathcal{X}_N$
 $x_0 = x$



Properties of fixed-complexity MPC:

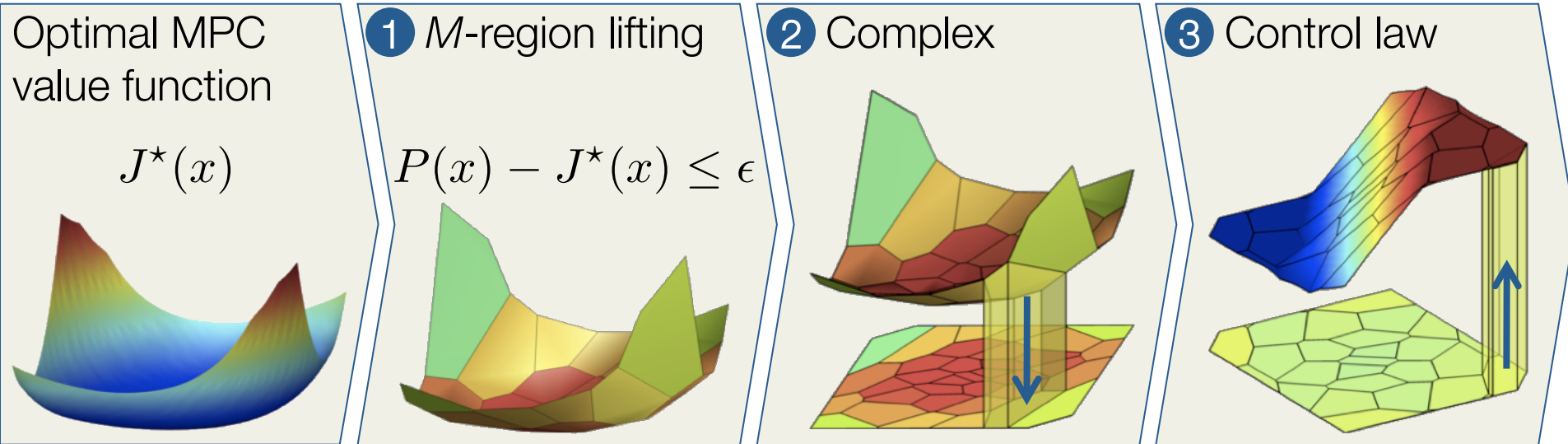
- Tradeoff between complexity and optimality
 - Real-time synthesis
 - Control extremely high-speed systems
- Process any convex MPC problem
- Synthesis of control law to software is verifiable

[Jones, Barić, Morari, 2007]

[Jones, Morari, 2010]

[Summers, Jones, Lygeros, Morari 2010]

Real-time explicit MPC : Offline processing



Given optimal controller:

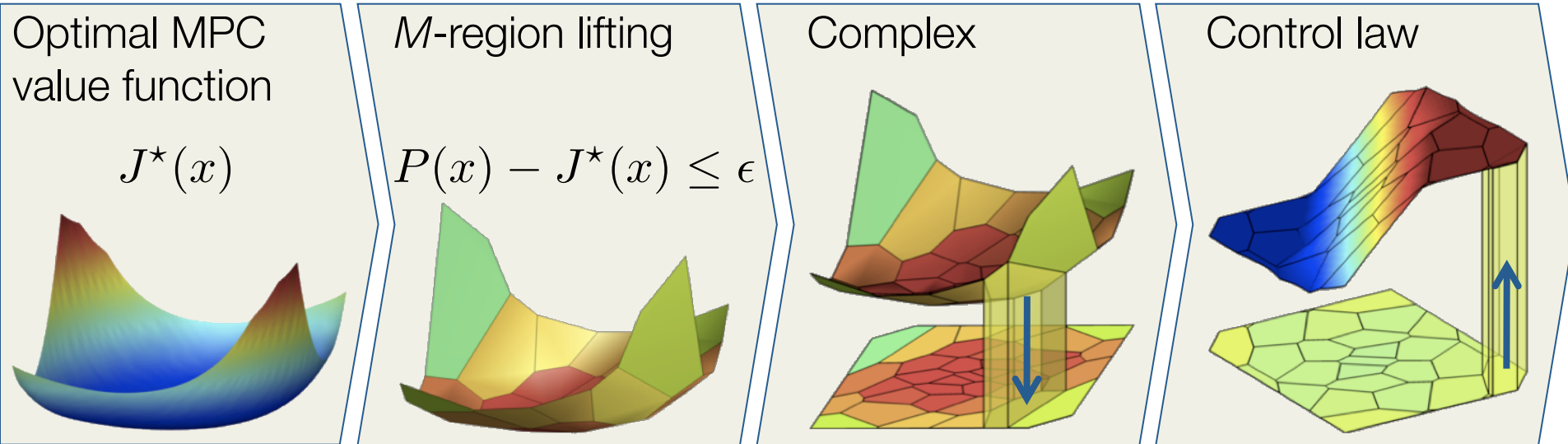
- 1 Compute convex polyhedral function of M facets
- 2 Define complex as projection of lifting facets
- 3 Interpolate optimal control law at vertices of complex

$$J^*(x_0) = \min_{u_i} V_N(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$

s.t. $x_{i+1} = f(x_i, u_i)$
 $(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$
 $x_N \in \mathcal{X}_N$

Result : Piecewise polynomial controller of M regions

Real-time explicit MPC : Properties



Real-time explicit MPC:

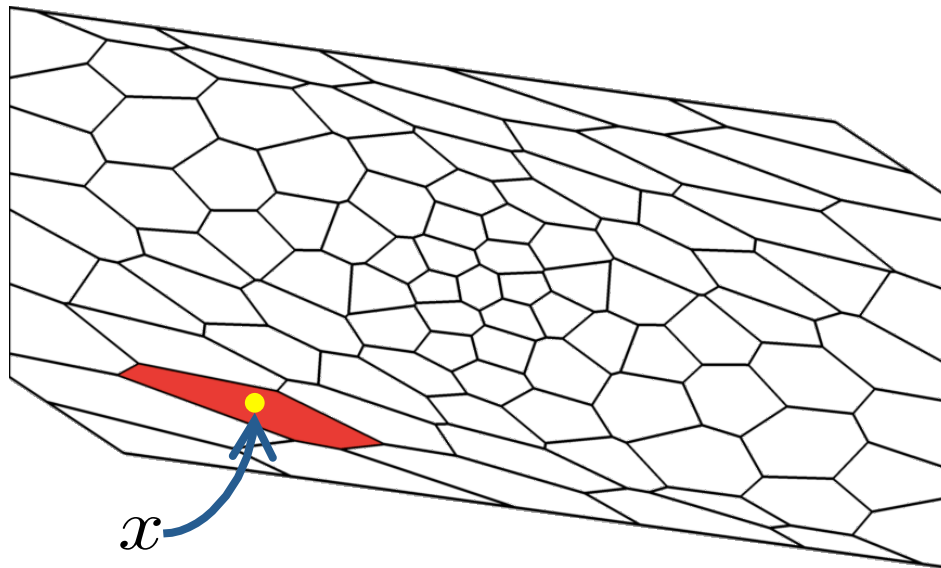
- Is computable in micro- to nanoseconds
- Satisfies constraints
- Stabilizes the system
- Complexity/performance tradeoff

Computational bottleneck : Point Location

Goal : General class of functions that can be evaluated in μs / ns

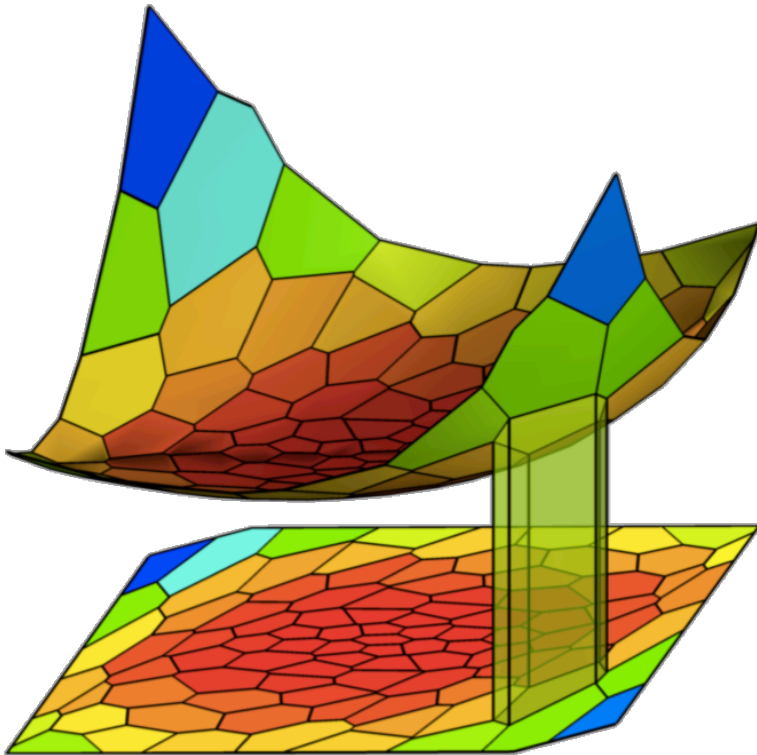
Critical operation : point location

Given $x \in \mathbb{R}^n$ and cell complex $\{C_1, \dots, C_m\}$ find i such that $x \in C_i$



Liftable complexes are log-time computable

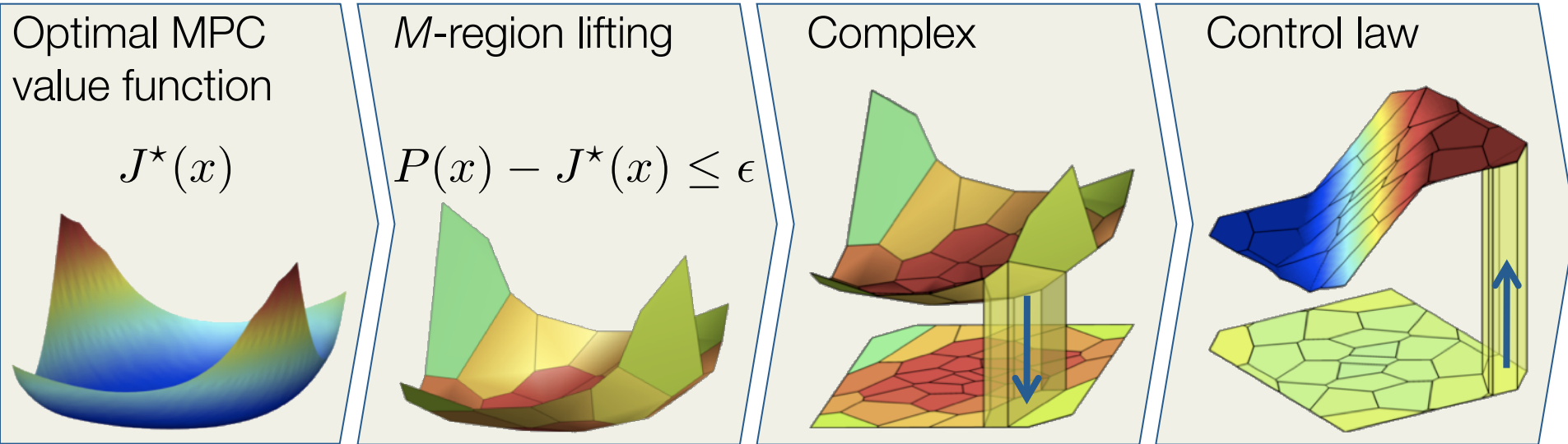
Goal : General class of functions that can be evaluated in μs / ns



- A polytopic complex is a power diagram iff it has a convex lifting
[Aurenhammer, 1991]
- Power diagram point location can be done in $O(\log n)$
[Mount et al, 1998]
- Explicit MPC with PWA cost has a lifting, but quadratic does not
[Jones et al, 2006]

Result : Design convex lifting \Rightarrow Log-time evaluation

Real-time explicit MPC : Properties



Real-time explicit MPC:

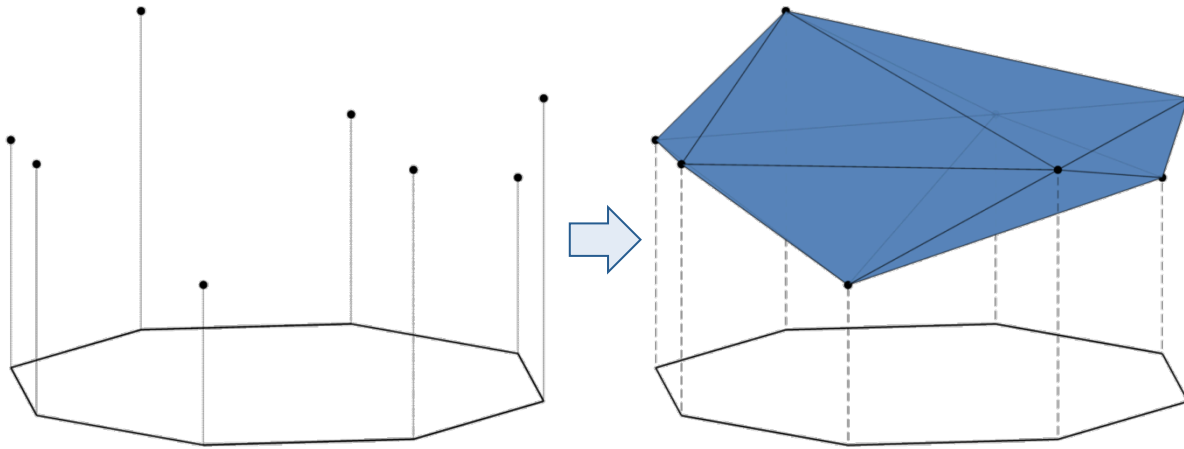
Is computable in micro- to nanoseconds \Leftarrow Lifting function

Satisfies constraints

Stabilizes the system

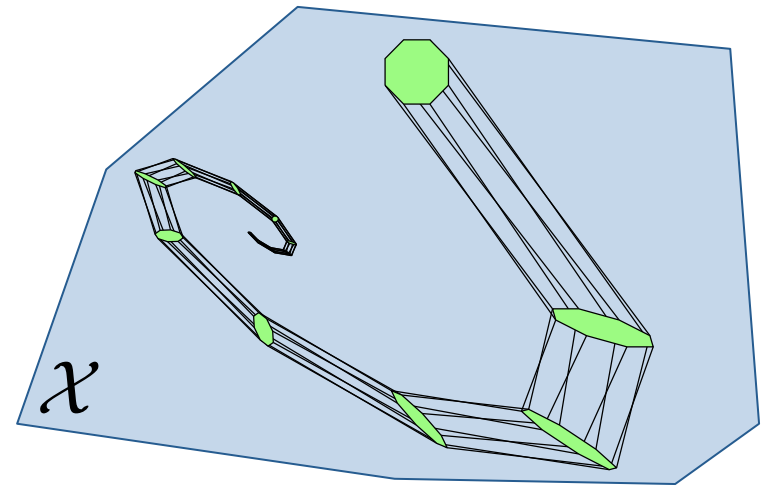
Complexity/performance tradeoff

Constraint satisfaction : Barycentric interpolation



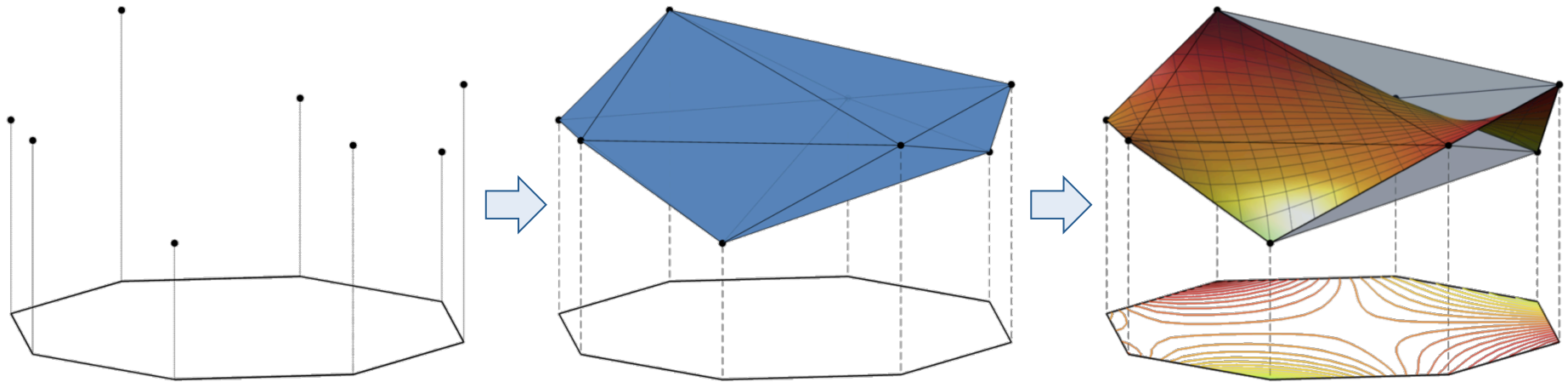
Goal : Interpolate vertices and satisfy constraints

- Convexity : Any interpolation inside convex hull is feasible



Result : Vertices feasible \Rightarrow Convex hull feasible

Constraint satisfaction : Barycentric interpolation



Goal : Interpolate vertices and satisfy constraints

- Convexity : Any interpolation inside convex hull is feasible

⇒ Barycentric interpolation

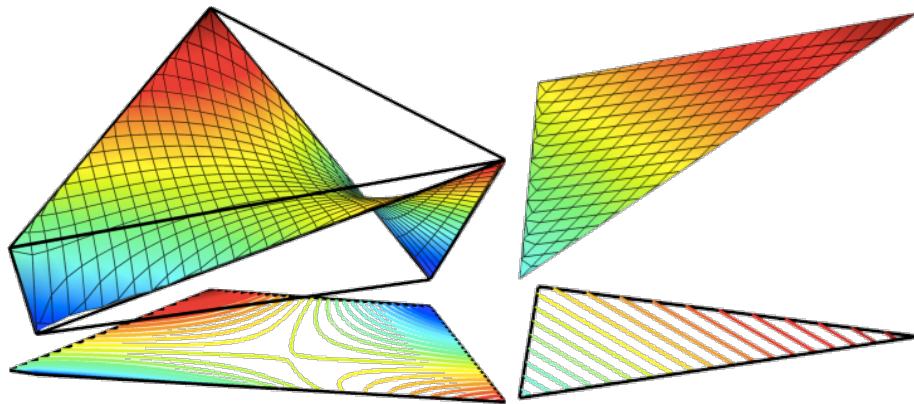
⇒ How to define w ?

$$\begin{aligned} w_v(x) &\geq 0 && \text{positivity} \\ \sum_{v \in \text{extreme}(S)} w_v(x) &= 1 && \text{partition of unity} \\ \sum_{v \in \text{extreme}(S)} v w_v(x) &= x && \text{linear precision} \end{aligned}$$

Constraint satisfaction : Barycentric interpolation

Thm: $\tilde{u}(x) = \sum_{v \in V} \frac{u^*(v) \alpha_v}{\|v - x\|_2}$
is barycentric for $\text{conv}(V)$

- α_v : area of facet individual polytope (pre-computed)
- Valid for *any polytope*
- Low data storage
- Evaluation in μs

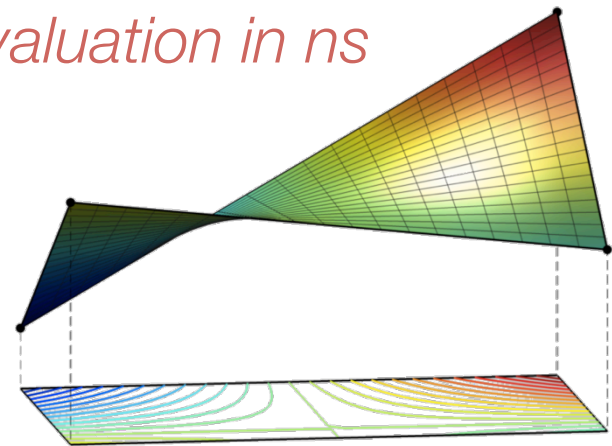


[Schaefer et al, 2008]

Thm: Tensor-product expansion of second-order interpolants is barycentric

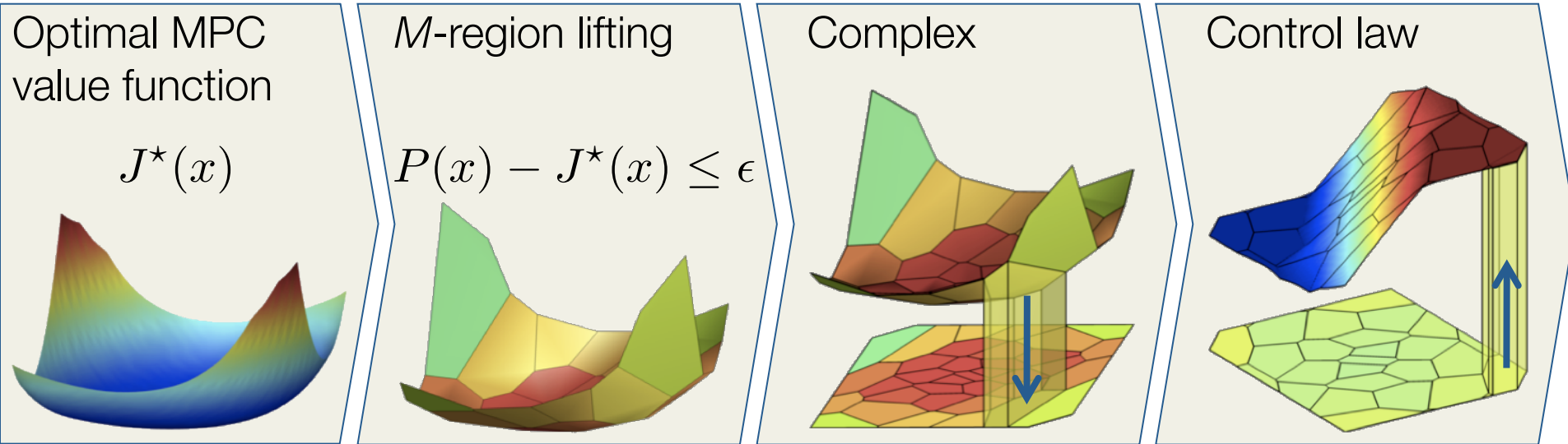
$$\tilde{u}(x) = \sum_v u^*(v) \prod_{j=1}^d \max \left\{ 0, \frac{|x_j - v_j| + 1}{h} \right\}$$

- Defined on hierarchical grid
- High data storage
- *Evaluation in ns*



[Summers, Jones, Lygeros, Morari 2009]

Real-time explicit MPC : Properties



Real-time explicit MPC:

Is computable in micro- to nanoseconds

<= Lifting function

Satisfies constraints

<= Barycentric interpolation

Stabilizes the system

Complexity/performance tradeoff

ϵ -approx controller is stable if $\epsilon < 1$

$$J(u) := V_N(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$

$$J^*(x_0) := \min_{u_i} J(u)$$

$$\text{s.t. } x_{i+1} = f(x_i, u_i)$$

$$(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$$

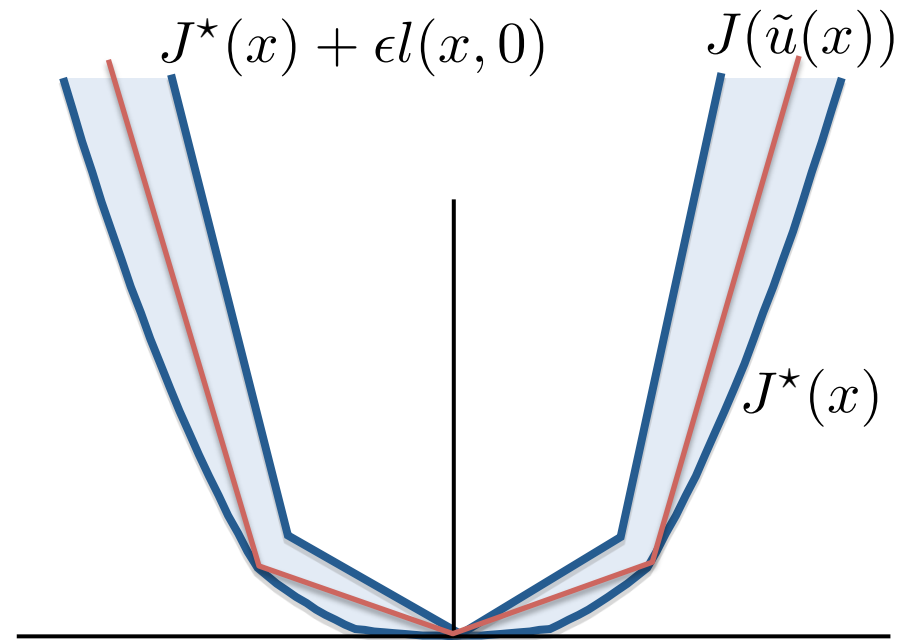
$$x_N \in \mathcal{X}_N$$

- Find a lifting use it to define $\tilde{u}(x)$
- Sufficiently close to optimal \Rightarrow Stabilizing

Thm: $x^+ = f(x, \tilde{u}(x))$
is stable if

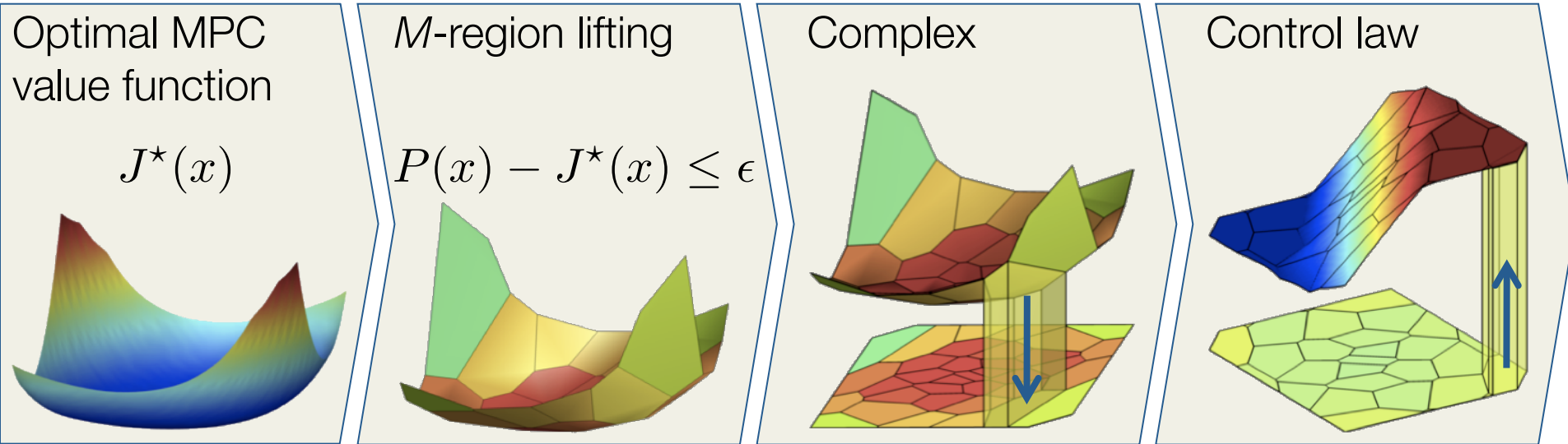
$$J^*(x) \leq J(\tilde{u}(x)) \leq J^*(x) + \epsilon l(x, 0)$$

for $\epsilon < 1$



Convex optimization verifies stability condition
without knowing optimal explicit solution

Real-time explicit MPC : Properties



Real-time explicit MPC:

Is computable in micro- to nanoseconds

Satisfies constraints

Stabilizes the system

\Leftarrow Lifting function

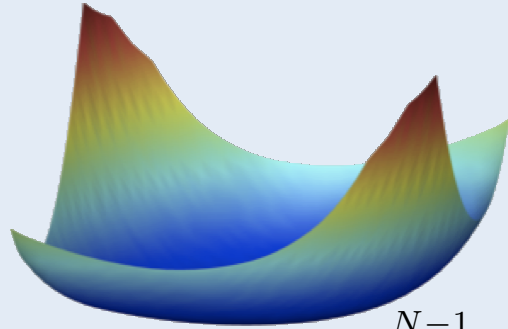
\Leftarrow Barycentric interpolation

\Leftarrow Error less than one

Complexity/performance tradeoff

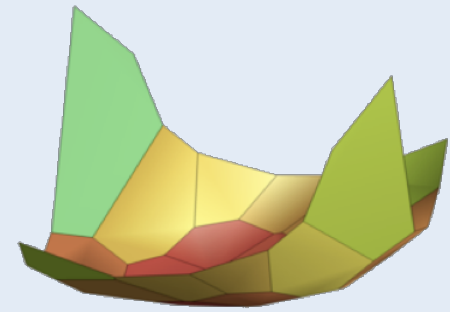
Approximate Convex Parametric Programming

M -region approximation \Rightarrow Double description method



$$J^*(x_0) = \min_{u_i} V_N(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$

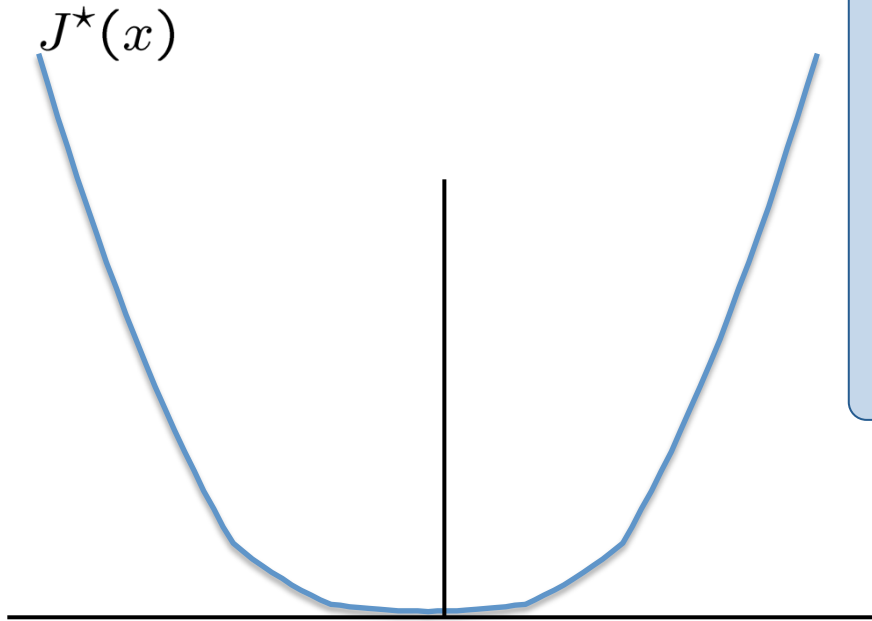
s.t. $x_{i+1} = f(x_i, u_i)$
 $(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$
 $x_N \in \mathcal{X}_N$



$$P(x)$$

- Open problem in many areas:
 - Vertex enumeration, Projection, Non-negative matrix factorization...
 - These problems are known to be NP-hard
 - Poly-time greedy-optimal algorithm
- \Rightarrow Lifting of M regions \Rightarrow Iterate algorithm M times!

Polyhedral approximation



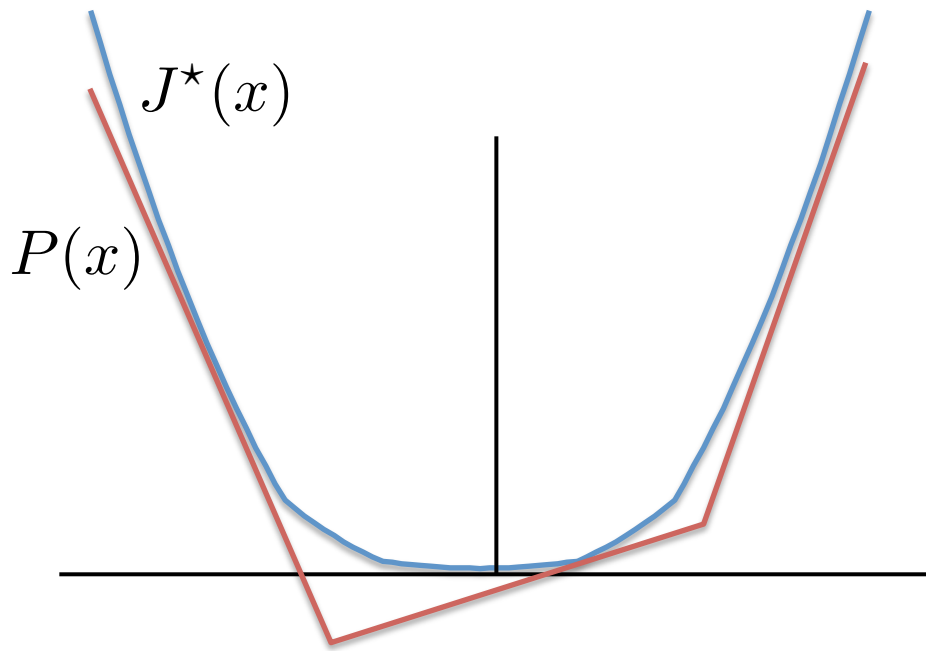
$$J^*(x_0) = \min_{u_i} V_N(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$
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$$(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$$
$$x_N \in \mathcal{X}_N$$

$\text{epi } J^*(x)$

- Implicitly defined
- Convex set
- Lyapunov function

Goal: Greedy-optimal polyhedral approx. of M -facets

Polyhedral approximation

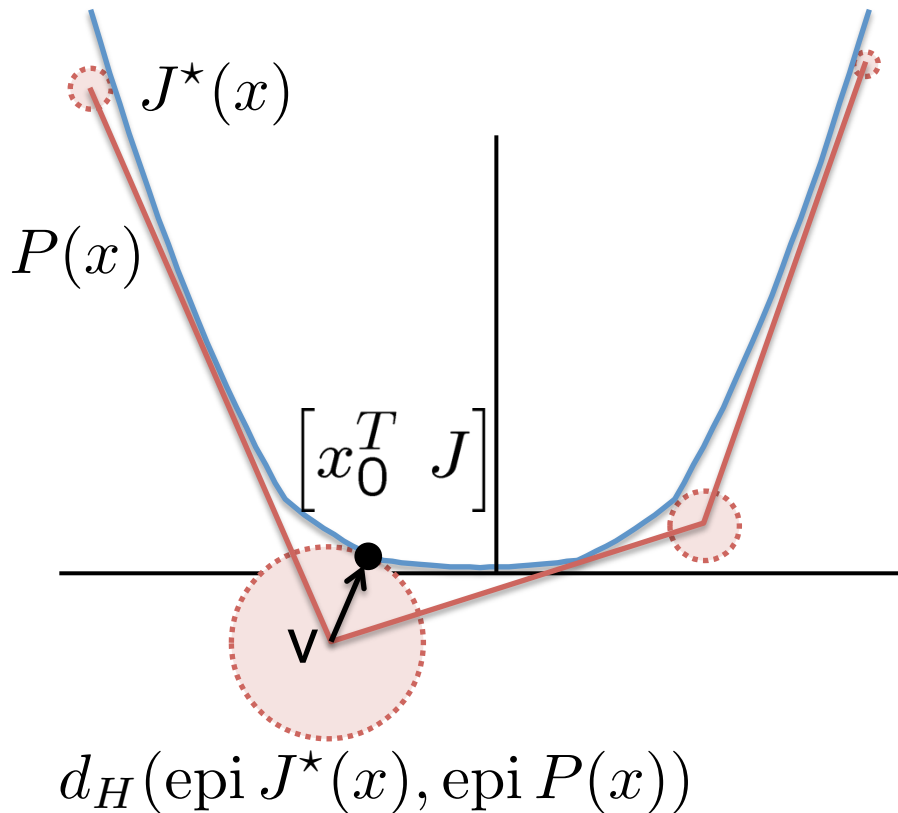


1. Outer polyhedral approx

Compute Hausdorff distance

$$\begin{aligned} & d_H(\text{epi } J^*(x), \text{epi } P(x)) \\ &= \max_{v \in \text{epi } P(x)} \min_{y \in \text{epi } J^*(x)} \|v - y\| \end{aligned}$$

Polyhedral approximation



1. Outer polyhedral approx
2. Project vertices onto optimal cost

$$\min \left\| v - \begin{bmatrix} x_0^T & J \end{bmatrix}^T \right\|_2^2$$

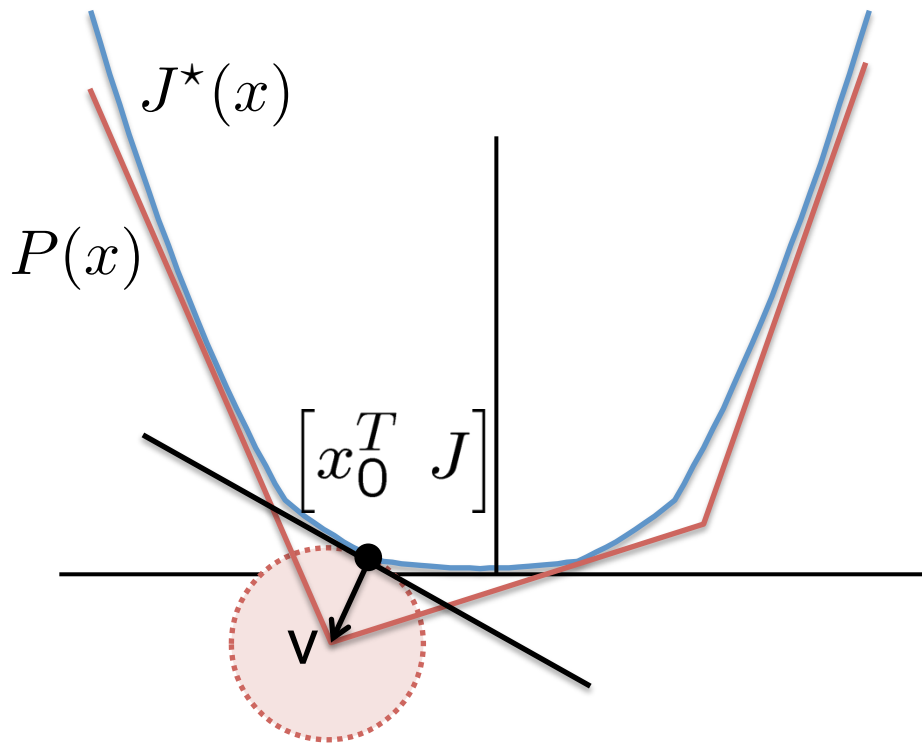
$$\text{s.t. } J \geq V_N(x_N) + \sum_{k=0}^{N-1} l(x_k, u_k)$$

$$x_{i+1} = f(x_i, u_i)$$

$$(x_i, u_i) \in \mathcal{X} \times \mathcal{U}$$

- Convex optimization!
- Do not need to know optimal solution

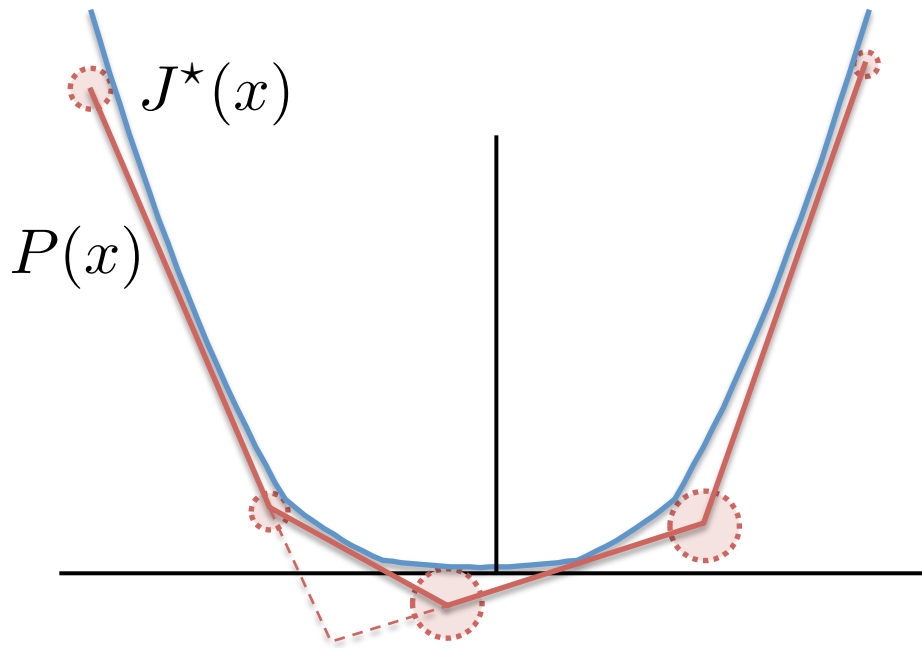
Polyhedral approximation



1. Outer polyhedral approx
2. Project vertices onto optimal cost
3. Insert maximally separating hyperplane

Maximally reduce Hausdorff distance

Polyhedral approximation

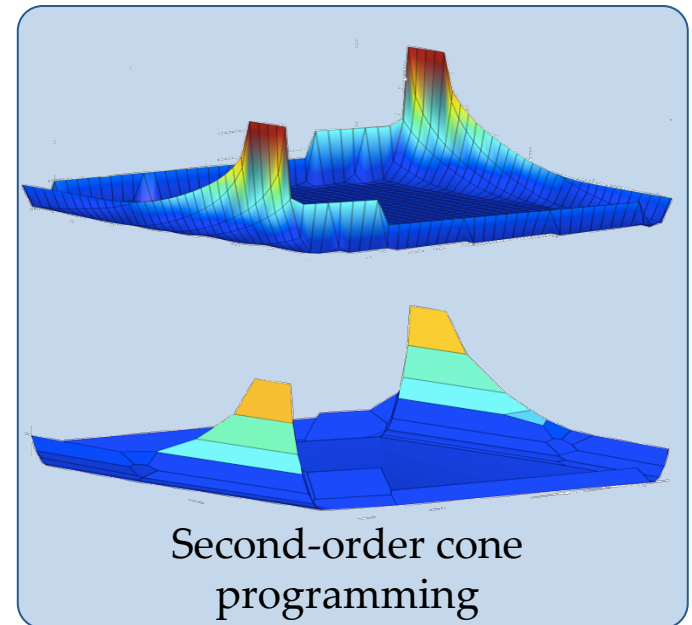
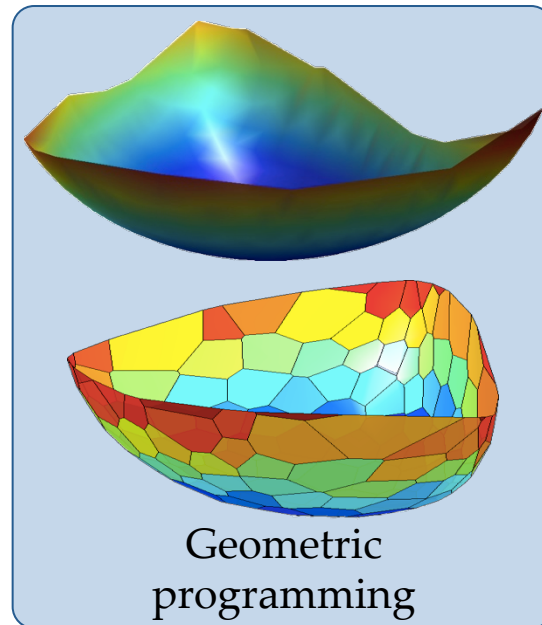
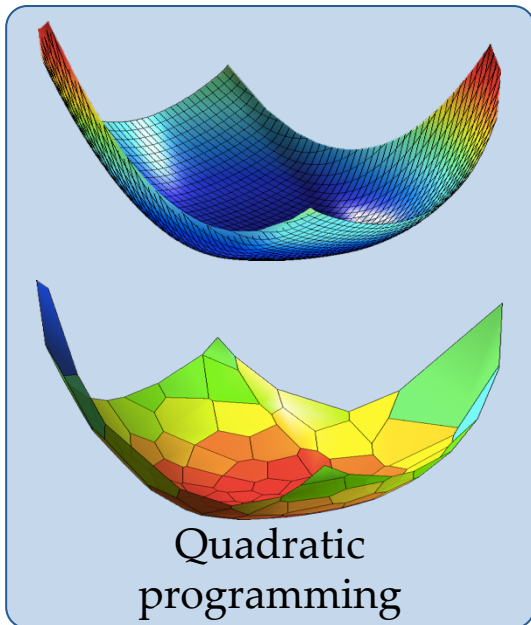


1. Outer polyhedral approx
2. Project vertices onto optimal cost
3. Insert maximally separating hyperplane
4. Repeat

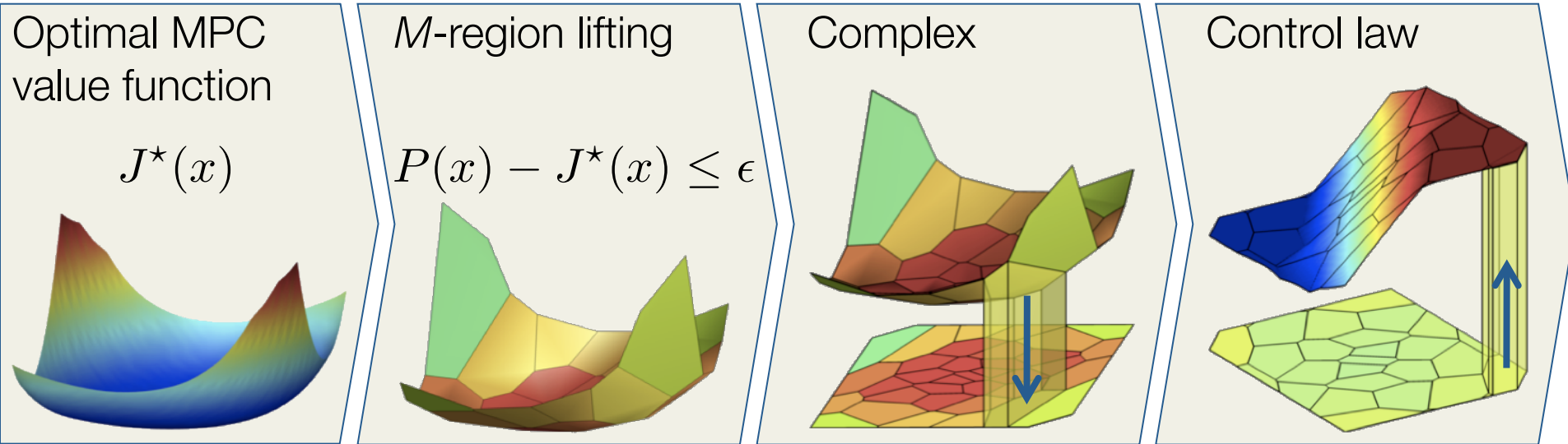
Result: M-region approximation with greedy-minimal Hausdorff distance.

Lifting calculation : Algorithm properties

- Lifting of M regions \leq Iterate algorithm M times
- Monotonic decrease in Hausdorff distance
 - Complexity / performance tradeoff via M
- There exists a minimum M for stability
 - ε -error in finite time \Rightarrow will find a Lyapunov function
 - Once stable, always stable



Real-time explicit MPC : Properties



Real-time explicit MPC:

Is computable in micro- to nanoseconds

Satisfies constraints

Stabilizes the system

Complexity/performance tradeoff

<= Lifting function

<= Barycentric interpolation

<= Error less than one

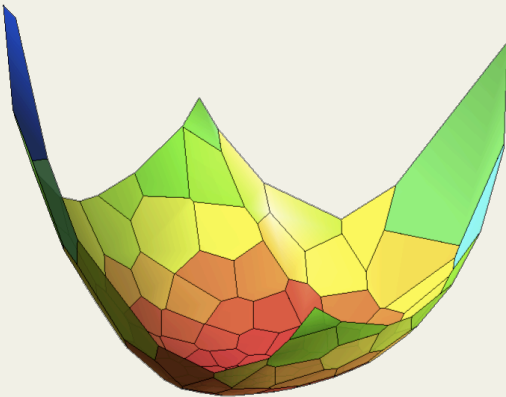
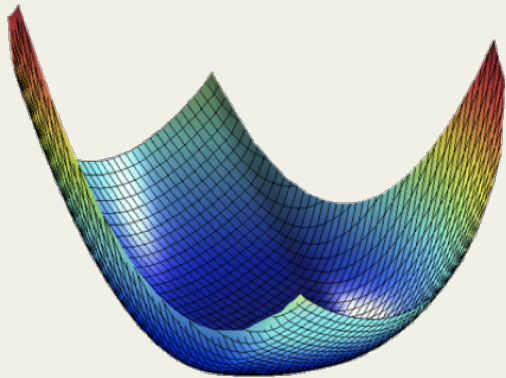
<= M-region lifting

Verified stability, feasibility & time without knowing optimal solution

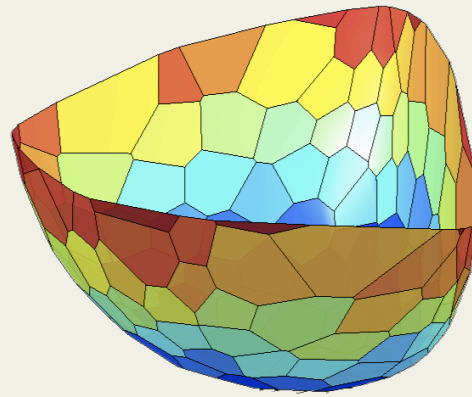
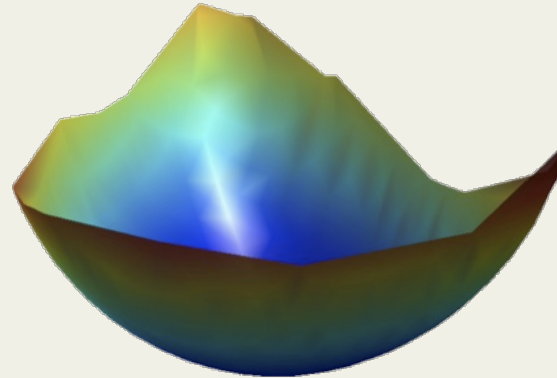
Double description method

Applicable to all convex parametric / set operations

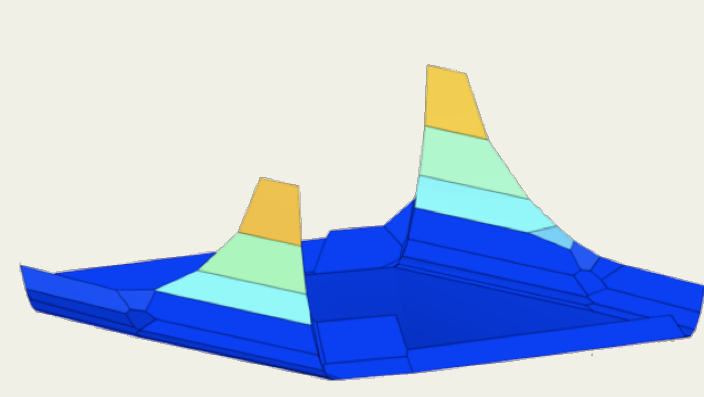
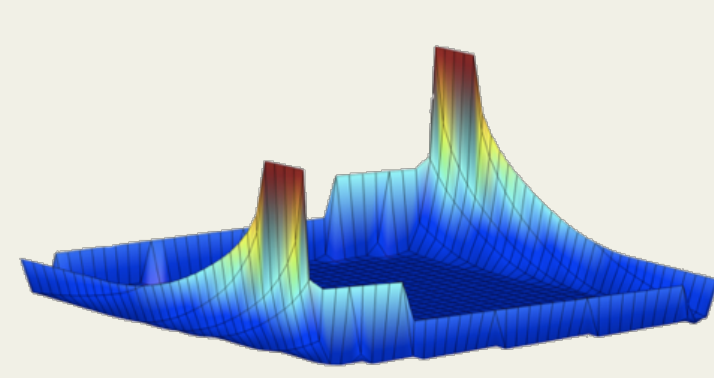
Quadratic programming



Geometric programming



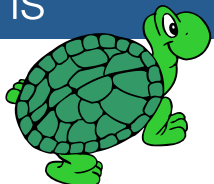
Second-order cone programming



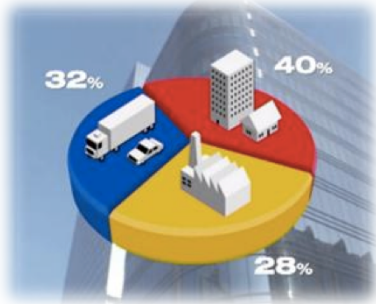
Applications by the Automatic Control Lab



18 ns	Multi-core thermal management (EPFL)	[Zanini et al 2010]
10 μ s	Voltage source inverters	[Mariethoz et al 2008]
20 μ s	DC/DC converters (STM)	[Mariethoz et al 2008]
25 μ s	Direct torque control (ABB)	[Papafotiou 2007]
50 μ s	AC / DC converters	[Richter et al 2010]
5 ms	Electronic throttle control (Ford)	[Vasak et al 2006]
20 ms	Traction control (Ford)	[Borrelli et al 2001]
40 ms	Micro-scale race cars	
50 ms	Autonomous vehicle steering (Ford)	[Besselmann et al 2008]
500 ms	Energy efficient building control (Siemens)	[Oldewurtel et al 2010]



Example: Energy Efficient Building Climate Control



Building sector: 40% of energy worldwide!
[International Energy Agency 2008]

Idea: Incorporate weather predictions into building controller

“The trouble with weather forecasting is that it's right too often for us to ignore it and wrong too often for us to rely on it.”

- Patrick Young

SIEMENS

Building Technologies

MeteoSwiss
Schweizerische Eidgenossenschaft
Confédération suisse
Confederazione Svizzera
Confederaziun svizra

EMPA
Materials Science & Technology

swiss**electric**
research

CCEM-CH
Competence Center Energy and Mobility

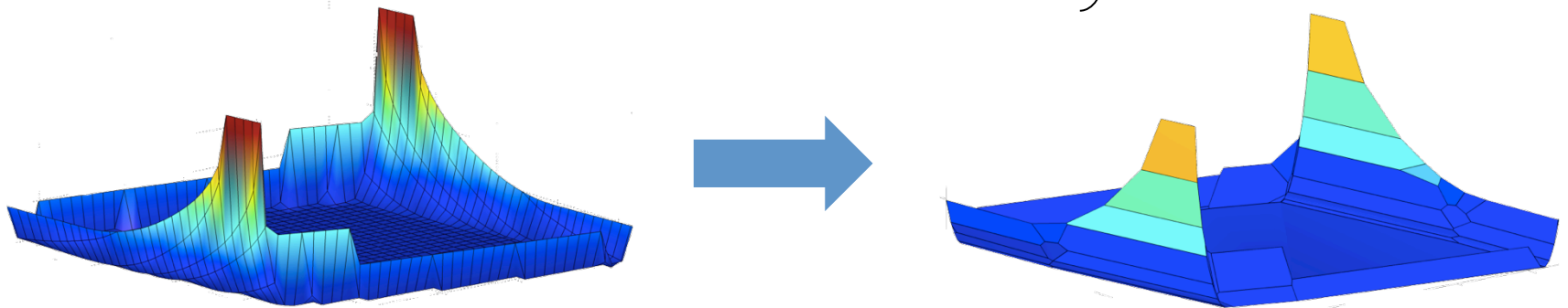
EPFL

AUTOMATIC CONTROL
LABORATORY

Example: Energy Efficient Building Climate Control

- Analysis shows uncertainty in weather forecast » Gaussian
- Stochastic model predictive controller
 - Constraints violated with a specified probability

$$\begin{aligned} \min \mathbb{E} & \left[\sum_{k=0}^{N-1} c^T \cdot \mu_k(\phi_k(x_0, \mu, \mathbf{w})) \right] \\ \text{s.t. } & \mu_k(\phi_k(x_0, \mu, \mathbf{w})) \in \mathcal{U} \quad \forall \mathbf{w} \in \mathbb{R}^{mN} \\ & \mathbb{P}\{\phi_k(x_0, \mu, \mathbf{w}) \in \mathcal{X}\} \geq 1 - \alpha \end{aligned} \quad \left. \vphantom{\begin{aligned} \min \mathbb{E} & \left[\sum_{k=0}^{N-1} c^T \cdot \mu_k(\phi_k(x_0, \mu, \mathbf{w})) \right] \\ \text{s.t. } & \mu_k(\phi_k(x_0, \mu, \mathbf{w})) \in \mathcal{U} \quad \forall \mathbf{w} \in \mathbb{R}^{mN} \\ & \mathbb{P}\{\phi_k(x_0, \mu, \mathbf{w}) \in \mathcal{X}\} \geq 1 - \alpha \end{aligned}} \right\} \begin{array}{l} \text{Convex} \\ \text{parametric} \\ \text{second-order} \\ \text{cone problem} \end{array}$$

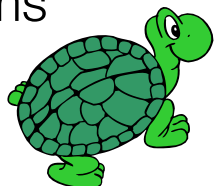
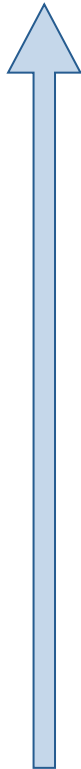


PWA controller with 30 regions => Can run in a light-switch
Energy savings in idealized Swiss buildings from 5% to 40%

Applications by the Automatic Control Lab



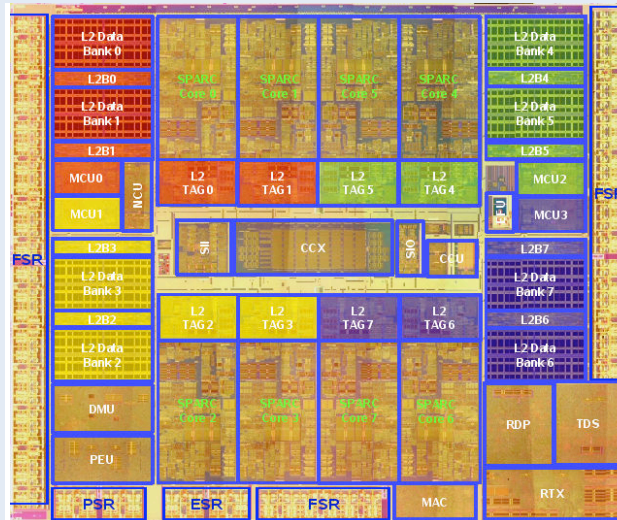
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Example :

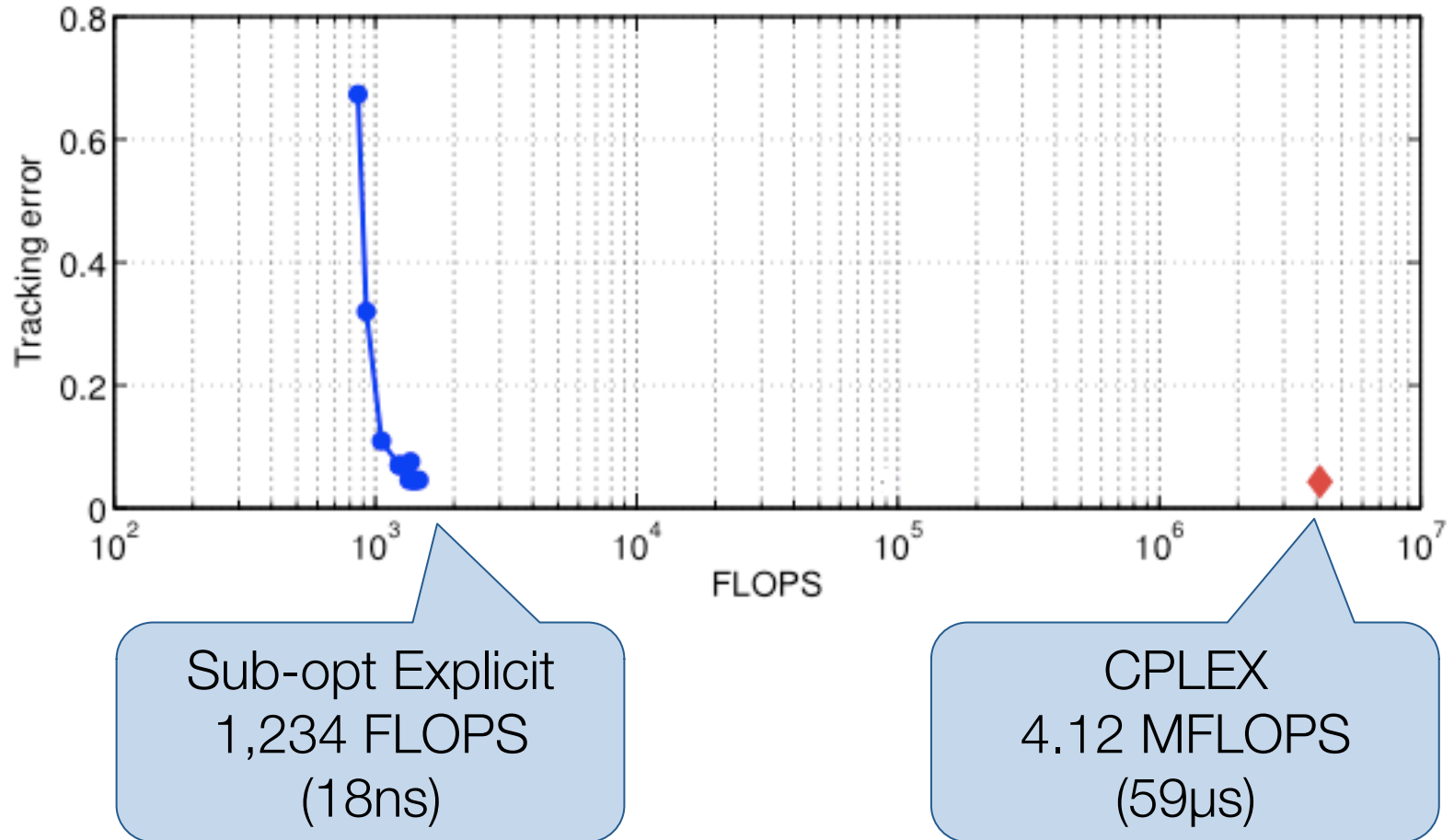
Temperature Regulation of Multi-Core Processor

- Goals
 - Track workload requests
 - Minimize power usage
 - Respect temperature limits
- Quadratic nonlinear dynamics
 - Exact convex relaxation
- Stringent computational and storage requirements



$$J^*(x_0, w) = \min_{f_i} \sum_{t=0}^N \sum_{i=0}^t (w_i - f_i)$$
$$\text{s.t. } x_{i+1} = Ax_i + Bf_i^2$$
$$\sum_{i=0}^t w_i \leq \sum_{i=0}^t f_i$$
$$x_i \leq T_{\max}$$
$$f_{\min} \leq f_i \leq f_{\max}$$

Computational results for QCQP : $>3,000\times$ faster



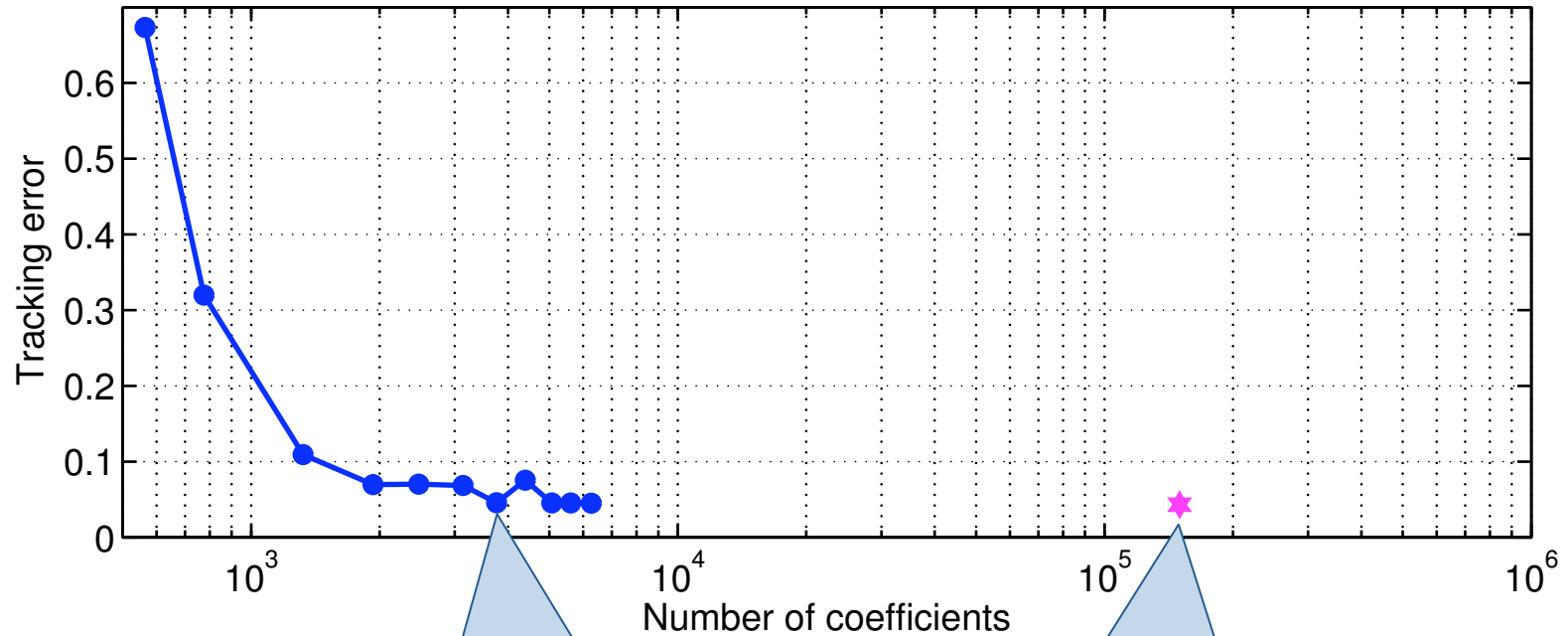
Sub-opt Explicit
1,234 FLOPS
(18ns)

CPLEX
4.12 MFLOPS
(59μs)

(Assuming 70 GFLOPS/sec – e.g., Intel Core i7 965 XE)

$>3,000\times$ faster than CPLEX

Computational results for QCQP : 45× less storage



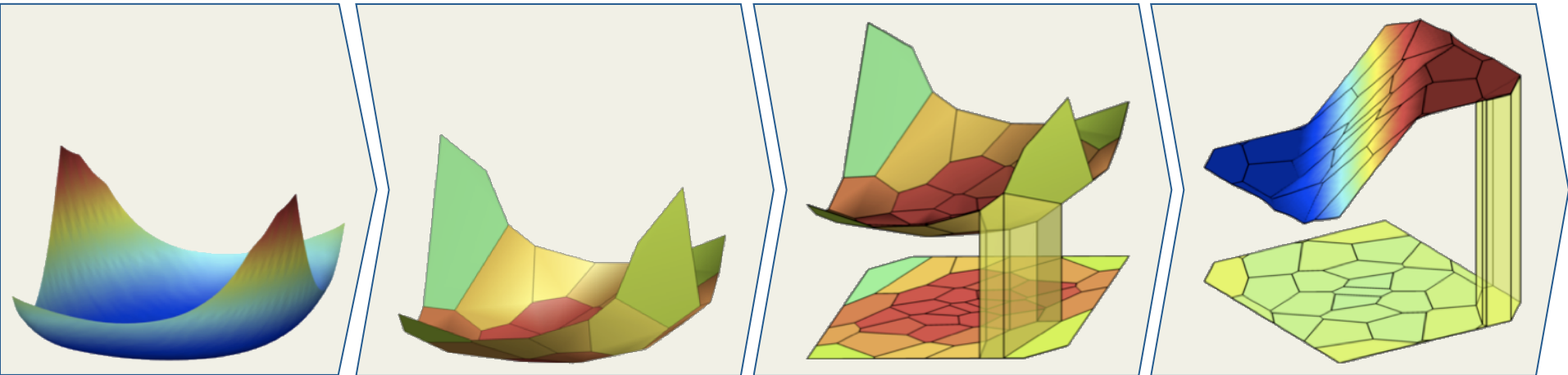
Sub-opt Explicit
26 KB

Linearised explicit
1.14 MB

45× less storage

Summary

- Complexity of explicit MPC control laws highly variable
 - Depends on tuning, dynamics, problem size, etc
- Fixed-complexity MPC
 - Produce sub-optimal control law of *specified* complexity
 - Certificate of stability, invariance
- Several related approximate explicit MPC methods
 - Different pros/cons
 - Based on sampling optimal control law and interpolating



Workshop Outline

14:00 – 14:30	Introduction
14:30 – 15:15	Lecture 1: MPC Theory
15:15 – 15:20	Break
15:20 – 15:40	Lecture 2 : Explicit MPC
15:40 – 16:00	Break
16:00 – 16:40	Lecture 2 : Explicit MPC
16:40 – 17:10	Lecture 3 : Real-time explicit MPC
17:10 – 17:20	Break
17:20 – 18:00	Lecture 4 : Real-time online optimization